

# Networking and Optimal Transport

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# Overview

## Overview of the presentation

- Optimal Transport on Network
- Leaf venation
- Interacting polymers
- The cavity approach

# General framework

- Distribution Networks:



$G = \{V, E\}$  : network

Supply of resource: Electricity, water, oxygen...

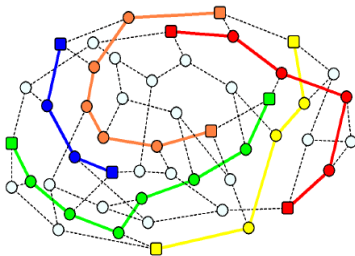
Nodes can be Sources and sinks. Find the configuration of currents that minimize the dissipation in the network

## Static Problems

# General Framework

- **Communication networks**

$a, b$  nodes; messages (...or car)  $\mu$  that go from  $a^\mu$  to  $b^\mu$ .



Find the paths that constitute the best compromise between total path length and network congestion

## Static Problems

# Distribution Network

- **Currents** are generated and absorbed on nodes ( $a = 1, \dots, N$ )

$i_a > 0$  Source ;

$i_a < 0$  Sink

- They circulate through edges:  
 $I_{ab}$  current circulating on edge ( $a \rightarrow b$ ) (directed)
- Current conservation (Kirchoff law).

$$i_a = \sum_b I_{ab}$$

# Distribution Network

- Energy dissipations  $J$

Edges; characterized by  
“distance”  $d_{ab}$  ; intrinsic  
conductance  $K_{ab} = k_{ab}d_{ab} = 1/R_{ab}$  ; they can be changed.

- $J = \sum_{(ab)} \frac{I_{ab}}{d_{ab}k_{ab}}$
- Cost to build up the network  $C[\{k_{ab}, d_{ab}\}]$

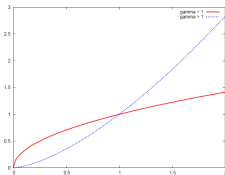
$$C[\{k_{ab}\}] = \sum_{ab} k_{ab}^{\gamma}$$

S. Bohn, M. Magnasco PRL 98, 088702 (2007)

# Distribution Network

- Minimize dissipation for fixed construction cost

$$C[\{k_{ab}\}] = \sum_{ab} k_{ab}^{\gamma} = K$$



- $\gamma > 1$  (Convex cost) many small transmission lines
- $\gamma < 1$  (Concave cost) few big transmission lines

# Result of optimization

- Triangular Lattice
- One source,  $N - 1$  sinks.
- Various values of  $\gamma$ .

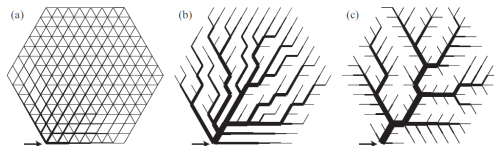


FIG. 2: Examples of the conductivity distributions. Results of the relaxation algorithm with different initial conditions: (a)  $\gamma = 2.0$  and (b)  $\gamma = 0.5$ . (c) Result with optimized topology with  $\gamma = 0.5$ . the arrow indicated the localized inlet, the remaining nodes are outlets with constant  $i_k$ .

- Unique (smooth) minimum for  $\gamma > 1$
- Local minima for  $\gamma < 1$ . Trees.



# A Little theorem

If  $C[\{k_{ab}\}]$  is concave the optimal transport network is a Tree.



Loops emerge from other needs:

- Robustness of transport to damage
- Fluctuations in the network use.

# Robustness of transport to damage

- Compute the power dissipated if  $P_{ab}$  if the edge  $(ab)$  is removed
- Consider as function to be optimized

$$R = \sum_{(ab)} P_{ab}$$

Katifori et al. PRL 2008

# Fluctuating Network Usage

- Assume that  $i_a$  are fluctuating with probability  $P(\mathbf{i})$
- Optimize the *average* dissipation

$$J_{av} = \sum_{ab} \frac{\langle I_{ab}^2 \rangle}{K_{ab}}$$

Katifori et al. PRL 2010

Carlson PRL 2010

# Loops & Hierarchy

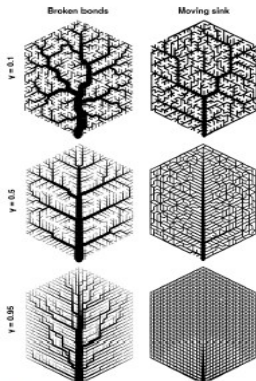
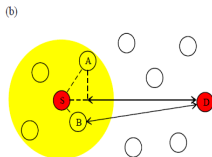
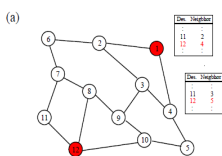


FIG. 2: Loops as a result of optimizing under damage to links (left column) and under a fluctuating load (right column). In all plots the vein thickness is proportional to  $C^{(2)/(1+(2)/\gamma)}$ .

# Communication Networks

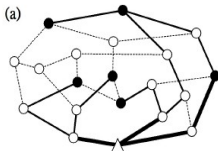
- **Networking:** routing, frequency allocation, information spreading, dynamic network allocation.
- **Routing:** Find the best route from  $a$  to  $b$ .



- Table based; Not dynamical ; insensitive to congestion
- Greedy: transmit to the neighbour close to destination

**Frequency allocation** to minimize interferences. Graph Coloring kind of problem.

# Routing and the Physics of Polymers



- Neglect time: Static Paths on networks.  $M$  paths  $M \sim N$ .  
Chi Ho Yeung and David Saad PRL 2012
- Avoid Congestion : Interaction between paths.

$$\sigma_{ab}^{\mu} = 0, \pm 1 \quad I_{ab} = \sum_{\mu} \sigma_{ab}^{\mu} - \sigma_{ba}^{\mu}$$

$$E[\{I_{ab}\}] = \sum_{ab} I_{ab}^{\alpha} \quad i_a^{\mu} = \sum_b \sigma_{ab}^{\mu}$$

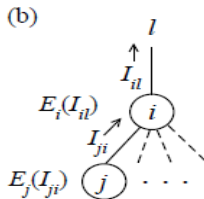
$\alpha > 1$ , minimize length, penalise congestion.

Polymers on networks with repulsive interaction.

At each node: A separate conservation law for each polymer.

# Cavity approach

- At each node: A separate conservation law for each polymer.
- Defining Messages on Edges.
- Message  $E_{a \rightarrow b}(\Sigma)$  ;  $\Sigma = \{\sigma^1, \dots, \sigma^M\}$



$$E_{a \rightarrow b}(\Sigma_{ab}) = I_{ab}^\alpha + \min_{\{I \mid \text{Constraints}\}} \sum_{c \in \partial a - b} E_{c \rightarrow a}(\Sigma_{ca})$$

$2^M$  messages at each link. Simplifications needed.

# Cavity approach

- Multiple sender ; Single receiver.
- A unique conservation law at each node.

$$i_a = \sum_b I_{a,b}$$

$$E_{a \rightarrow b}(I_{ab}) = I_{ab}^\alpha + \min_{\{I \mid \text{Constraint}\}} \sum_{c \in \partial a - b} E_{c \rightarrow a}(I_{ca})$$

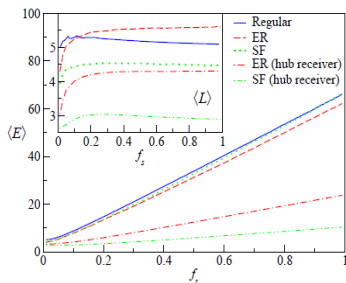
$|I| < M$ .  $\alpha > 1$  iterative procedure to compute the messages.  
Convex cost function.



## Some results

Study of average Energy diss. and Length as a function of the fraction of senders  $f$  for different graph models.

- Erdos-Renyi, Random regular, Scale free



# Some generalizations

- Non Convex cost functions
- Different kind of Receiver and Senders
- Effect of exclusion
- Free-flow and Jammed Phases.