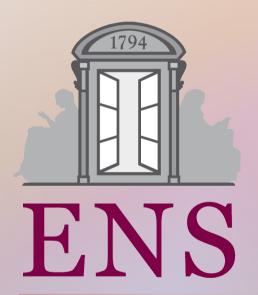
Alberto Guggiola

École Normale Supérieure - Paris

Epidemic processes on networks, viral marketing and optimal vaccination

with G. Semerjian







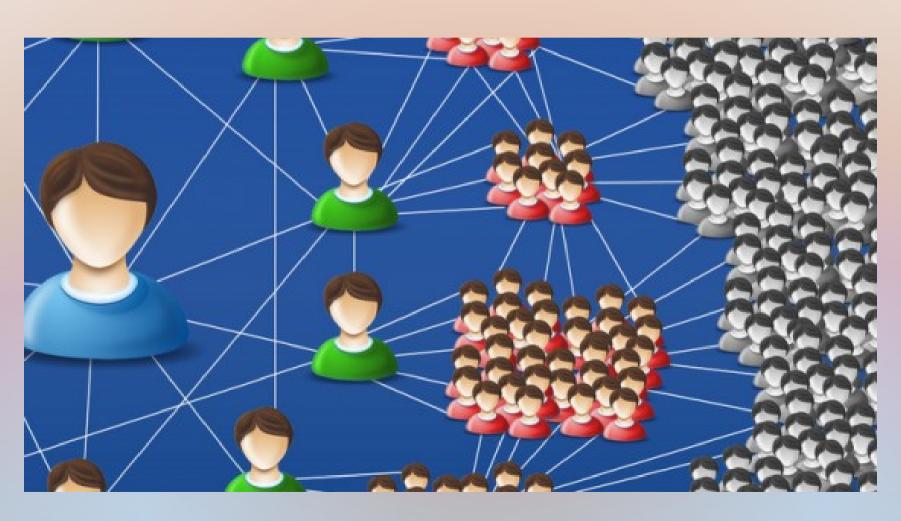
Statistical Physics Approaches Networks Across Disciplines



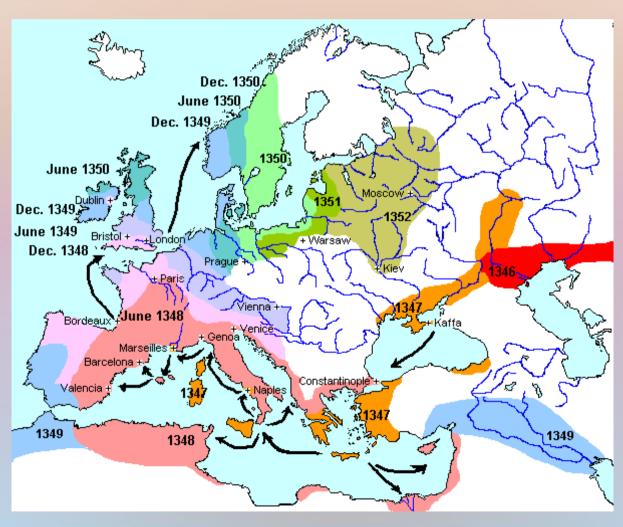
Cascade failures in power grids...



... effectiveness of viral marketing campaigns...



... diffusion of epidemics...



... shocks in financial markets...



Fundamental characteristics:

System:

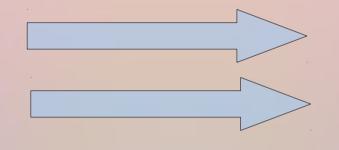
Elements
Interconnections

Fundamental characteristics:

System:

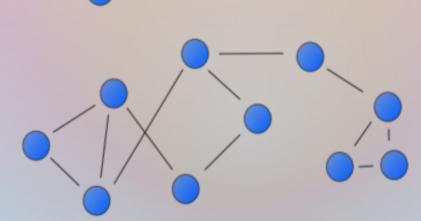
Network:

Elements
Interconnections



Nodes

Links



Individual

The Susceptible – Infective Model

• Two discrete states (Compartmental Model)

 Microscopically irreversible dynamics (no recovery)



The Dynamics

A node can be activated

From the beginning: Seed of the infection At time t:
At least 1 active neighbors
at time t - 1

D Kempe, J M Kleinberg, E Tardos, Maximizing the spread of influence through a social network, KDD 2003.

The Dynamics

A node can be activated



From the beginning:

Seed of the infection

At time t:

At least 1 active neighbors at time t - 1

$$\Psi_i = \mathbf{1}[t_i = 0] + \mathbf{1}[t_i = f(\{t_j\})] = 1$$

D Kempe, J M Kleinberg, E Tardos, Maximizing the spread of influence through a social network, KDD 2003.

Formalizing the objective

Different questions...

- How many infected at most with n seeds?
- How many seeds at least to activate everyone?
- Is everyone activated after enough time?
- If not, what percentage is not touched?

Formalizing the objective

... (almost) the same answer:

Minimization of:
$$E(t) = \sum E_i(t_i)$$

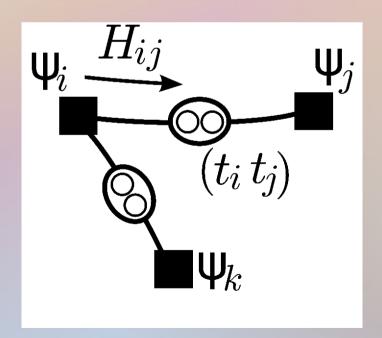
where
$$E_i(t_i) = \mu \mathbf{1}[t_i = 0] + \varepsilon \mathbf{1}[t_i = \infty]$$

F Altarelli, A Braunstein, L Dall'Asta, R Zecchina, Optimizing spread dynamics on graphs by message passing, J Stat Mech (2013)

The Message Passing Algorithm

Efficient representation of the variables and the constraints of the problem

Constraints: $\Psi_i = 1$, $\forall i$ Variables: (t_i, t_j) If nodes i and j are linked together



Consistent activation times in absence of a neighbor $H_{ii}(t_i, t_i)$ are searched iteratively

As the procedure converges, the distribution of the true activation times can be read

Two Simpler Algorithms

Random Choice

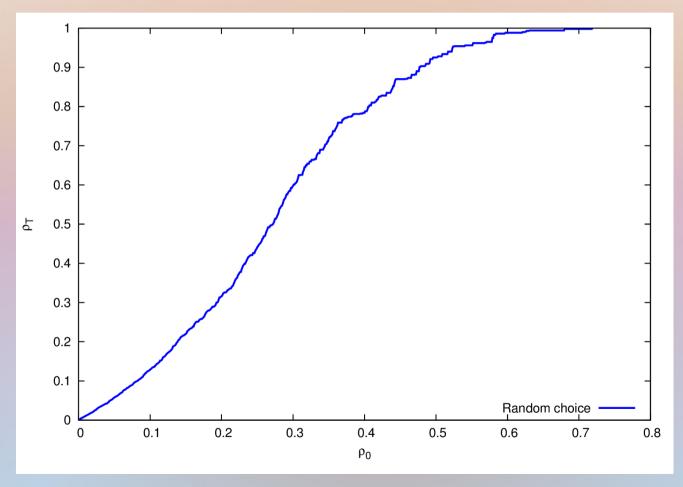
New randomly chosen seeds are added until a configuration activating all the network is reached.

Greedy Strategy

Each extra seed is chosen so to increment the most the number of activated nodes w.r.t. the previous seed configuration.

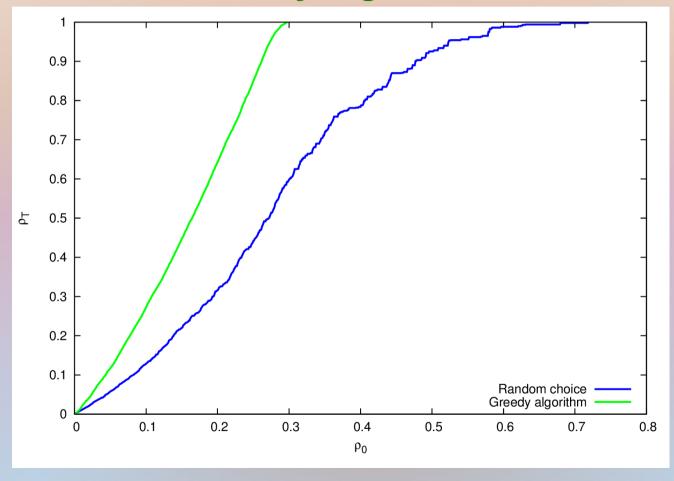
N=1000, k=3, l=2, T=5

Random choice



N=1000, k=3, l=2, T=5

Random choice Greedy algorithm

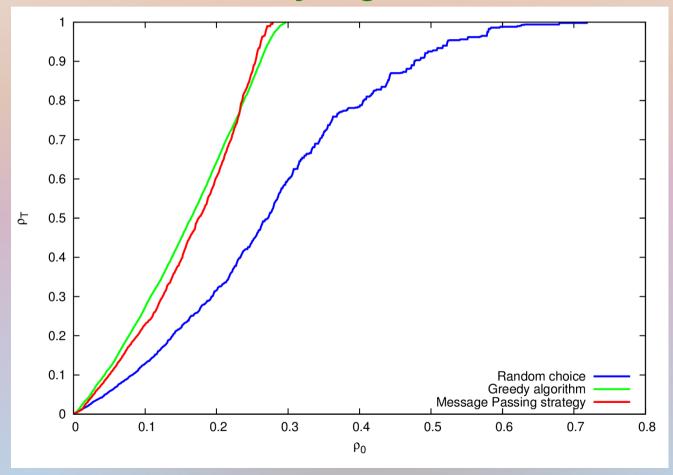


N=1000, k=3, l=2, T=5

Random choice

Greedy algorithm

MP strategy*



* 1 step RSB

Some analytical results

Seeds needed for complete activation in an infinite time: Random vs optimal choice

$$\theta_{optimal} = \frac{k-2}{2(k-1)} = \frac{1}{2} \theta_{random}$$

For a k – Random Regular Graph

$$l = k-1; T \rightarrow \infty; N \rightarrow \infty$$

Perspectives

Application to different types of graphs

Application to real world networks (e.g. biological, economical, infrastructural)

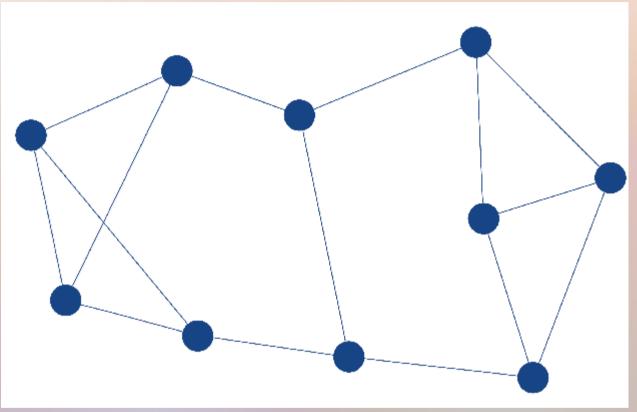
Perspectives

Application to different types of graphs

Application to real world networks (e.g. biological, economical, infrastructural)

Thank you for your attention!

Random Regular Graphs



Random Regular Graphs

Each node has k neighbours, randomly chosen

RRG with N = 10 nodes and k = 3

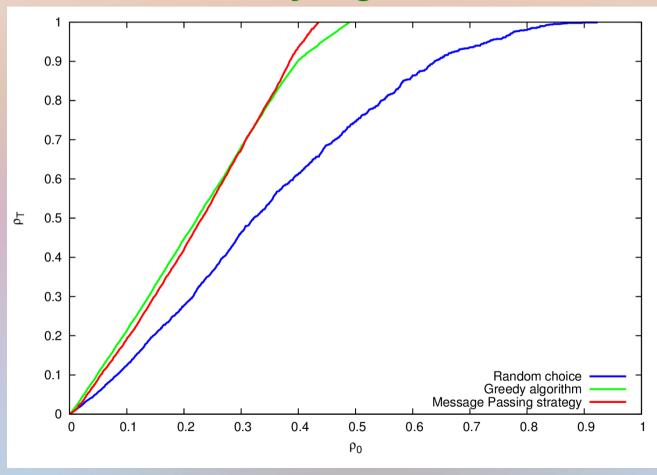
R Albert, AL Barabàsi, Statistical Mechanics of Complex Networks, Rev. Mod. Phys. 74, 47–97 (2002)

N=1000, k=3, l=2, T=1

Random choice

Greedy algorithm

MP strategy

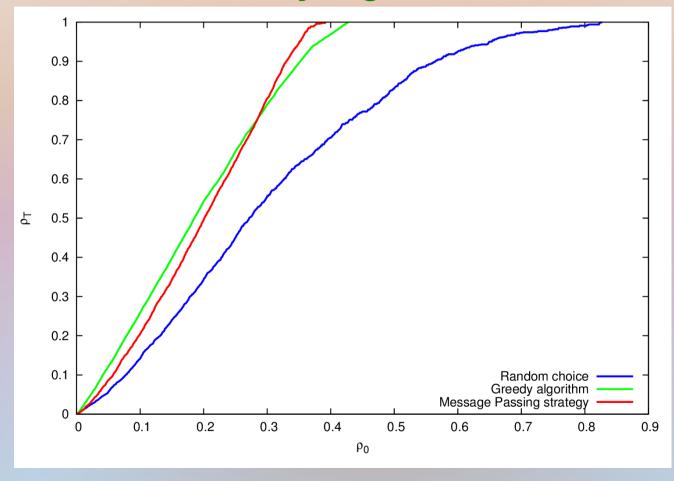


N=1000, k=4, l=2, T=1

Random choice

Greedy algorithm

MP strategy



N=800, k=3, l=2, T=5

Random choice

Greedy algorithm

MP strategy

